

Generic Synchronization Policies in C++

Ciaran McHale

CiaranMcHale.com
Complexity explained simply

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Introduction

- Most people know that writing synchronization code is:
 - Difficult: APIs are low-level
 - Non-portable: many threading APIs: POSIX, Windows, Solaris, DCE, ...
- In practice, most synchronization code implement a small number of high-level “usage patterns”:
 - Let's call these *generic synchronization policies* (GSPs)
 - The most common GSPs can be implemented as a C++ library
- Using GSPs in applications:
 - Is much easier than using low-level APIs
 - Encapsulates the underlying threading package → provides portability

1. Scoped Locks

Critical section

- The following (pseudocode) function uses a critical section:

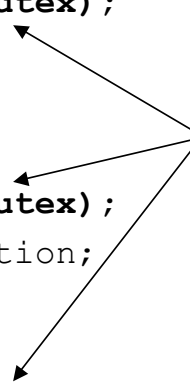
```
void foo()  
{  
    getLock (mutex) ;  
    ...  
    releaseLock (mutex) ;  
}
```

- The above code is very simple. However...
- Complexity increases if the function has several exit points:
 - Because `releaseLock()` must be called at each exit point
 - Examples of extra exit points:
 - Conditional `return` statements
 - Conditionally throwing an exception

Critical section with multiple exit points

```
void foo()  
{  
    getLock (mutex) ;  
    ...  
    if (...) {  
        releaseLock (mutex) ;  
        return;  
    }  
    if (...) {  
        releaseLock (mutex) ;  
        throw anException;  
    }  
    ...  
    releaseLock (mutex) ;  
}
```

Have to call `releaseLock()` at every exit point from the function



Critique

- **Needing to call `releaseLock()` at every exit point:**
 - Clutters up the “business logic” code with synchronization code
 - This clutter makes code harder to read and maintain
 - **Forgetting to call `releaseLock()` at an exit point is a common source of bugs**
-
- **There is a better way...**

Solution: `ScopedMutex` class

- **Define a class called, say, `ScopedMutex`:**
 - This class has no operations! Just a constructor and destructor
 - Constructor calls `getLock()`
 - Destructor calls `releaseLock()`
- **Declare a `ScopedMutex` variable local to a function**
 - At entry to function → constructor is called → calls `getLock()`
 - At exit from function → destructor is called → calls `releaseLock()`
- **The following two slides show:**
 - Pseudocode implementation of `ScopedMutex` class
 - Use of `ScopedMutex` in a function

The ScopedMutex class

```
class ScopedMutex {
public:
    ScopedMutex(Mutex & mutex)
        : m_mutex(mutex)
        { getLock(m_mutex); }

    ~ScopedMutex()
        { releaseLock(m_mutex); }
private:
    Mutex & m_mutex;
};
```

Use of ScopedMutex

```
void foo()
{
    ScopedMutex    scopedLock(mutex);
    ...
    if (...) { return; }
    if (...) { throw anException; }
    ...
}
```

No need to call `releaseLock()` at every exit point from the function!

Comments on ScopedMutex

- This technique is *partially* well known in the C++ community:
 - 50% of developers the author worked with already knew this technique
 - They considered it to be a “basic” C++ coding idiom
 - Other 50% of developers had not seen the technique before
- Of the developers who already knew this technique:
 - They all used it for mutex locks
 - Only a few knew it could be used for readers-writer locks too
 - Nobody knew it could be used for almost any type of synchronization code
- Contribution of this presentation:
 - Generalize the technique so it can be used much more widely
- To explain how to do this, I need to take a slight detour:
 - Have to introduce the concept of *generic synchronization policies*

2. The Concept of Generic Synchronization Policies

Genericity for types

- C++ provides template types
- Example of a template type definition:

```
template<t> class List { ... };
```

- Examples of template type instantiation:

```
List<int>      myIntList;  
List<double>  myDoubleList;  
List<Widget> myWidgetList;
```

- Some other languages provide a similar capability, often with different terminology and syntax
 - Perhaps called *generic types* instead of *template types*
 - Perhaps surround type parameters with [] instead of <>

Genericity for synchronization policies

- Using a pseudocode notation, here are declarations of mutual exclusion and readers-writer policies

```
Mutex[Op]  
RW[ReadOp, WriteOp]
```

- In above examples, each parameter is a set of operations
- Example instantiations on operations Op1, Op2 and Op3

```
Mutex[{Op1, Op2, Op3}]  
RW[{Op1, Op2}, {Op3}]
```

Producer-consumer policy

- Useful when:

- A buffer is used to transfer data between threads
- A producer thread *puts* items into the buffer
- A consumer thread *gets* items from the buffer
- If the buffer is empty when the consumer tries to get an item then the consumer thread blocks
- The buffer might have *other* operations that examine the state of the buffer

- In pseudocode notation, the policy declaration is:

```
ProdCons[PutOp, GetOp, OtherOp]
```

- Example instantiations:

```
ProdCons[{insert}, {remove}, {count}]
ProdCons[{insert}, {remove}, {}]
```

Bounded producer-consumer policy

- Variation of the producer-consumer policy:

- Buffer has a fixed size
- If the buffer is full when the producer tries to put in an item then the producer thread blocks

- In pseudocode notation, policy is:

```
BoundedProdCons(int size)[PutOp, GetOp, OtherOp]
```

- Typically, the `size` parameter is instantiated on a parameter to the constructor of the buffer class

- An example instantiation will be shown later

3. Generic Synchronization Policies in C++

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Mapping Mutex[Op] into C++

```

class GSP_Mutex {
public:
    GSP_Mutex() { /* initialize m_mutex */ }
    ~GSP_Mutex() { /* destroy m_mutex */ }
    class Op {
    public:
        Op(GSP_Mutex & data) : m_data(data)
        { getLock(m_data.m_mutex); }
        ~Op()
        { releaseLock(m_data.m_mutex); }
    private:
        GSP_Mutex & m_data;
    };
private:
    friend class ::GSP_Mutex::Op;
    OS-specific-type m_mutex;
};

```

Class name = "GSP_"
+ name of policy

Constructor & destructor
of outer class initialize and
destroy locks

A nested class
for each policy
parameter

Constructor & destructor
of nested class *get* and
release locks stored in
the outer class

Mapping RW[ReadOp, WriteOp] into C++

```
class GSP_RW {
public:
    GSP_RW();
    ~GSP_RW();
```

```
class ReadOp {
public:
    ReadOp(GSP_RW & data);
    ~ReadOp();
};
```

```
class WriteOp {
public:
    WriteOp(GSP_RW & data);
    ~WriteOp();
};
```

```
};
```

This policy has two parameters so there are two nested classes

Mapping BoundedProdCons into C++

- This is the mapping for

```
BoundedProdCons(int size) [PutOp, GetOp, OtherOp]
```

```
class GSP_BoundedProdCons {
public:
    GSP_BoundedProdCons(int size);
    ~GSP_BoundedProdCons();
    class PutOp    {...};
    class GetOp    {...};
    class OtherOp  {...};
};
```

The size parameter to the policy maps into a parameter to the constructor of the class

This policy has three parameters so there are three nested classes

Instantiating GSP_RW[ReadOp, WriteOp]

```

#include "gsp_rw.h"

class Foo {
private:
    GSP_RW    m_sync;
public:

    void op1(...) {
        GSP_RW::ReadOp    scopedLock(m_sync);
        ...
    }

    void op2(...) {
        GSP_RW::WriteOp   scopedLock(m_sync);
        ...
    }
};

```

#include header file (name of class written in lowercase)

Add instance variable whose type is name of policy's outer class

Synchronize an operation by adding a local variable whose type is a nested class of the policy

Instantiating GSP_BoundedProdCons

```

#include "gsp_boundedprodcons.h"

class Buffer {
private:
    GSP_BoundedProdCons    m_sync;
public:
    Buffer(int size) : m_sync(size) { ... }

    void insert(...) {
        GSP_BoundedProdCons::PutOp    scopedLock(m_sync);
        ...
    }

    void remove(...) {
        GSP_BoundedProdCons::GetOp    scopedLock(m_sync);
        ...
    }
};

```

The size parameter of the policy is initialized with value of a parameter to the constructor

4. Critique

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Strengths of GSPs

- Only one person needs to know how to implement GSPs
 - Trivial for everyone else to instantiate GSPs
- Separates synchronization code from “business logic” code
 - Improve readability and maintainability of both types of code
- Removes a common source of bugs:
 - Locks are released even if an operation throws an exception
- Improves portability:
 - API of GSPs does *not* expose OS-specific details of synchronization
- Efficiency:
 - GSPs can be implemented with inline code

Potential criticisms fo GSPs

- “Can they handle *all* my synchronization needs?”
 - 80/20 principle: *most* synchronization needs can be handled by just a small library of GSPs
 - You are not restricted to a library of pre-written GSPs. Instead...
 - You can write new GSPs if the need arises
- “GSPs are just a `ScopedMutex` with a new name”
 - The “just” part is inaccurate
 - GSPs generalize the `ScopedMutex` concept so it can be used for a much wider set of synchronization policies

Issues not addressed

- GSPs do not address:
 - POSIX thread cancellation
 - Timeouts
 - Lock hierarchies
- In the author’s work, these issues arise infrequently so he did not bother to support them
 - GSPs could probably be extended to support the above issues

5. Ready-to-run GSPs

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Ready-to-run GSPs

- A library of ready-to-use GSPs is available:
 - Download from www.CiaranMcHale.com/download
 - Documentation provided in multiple formats:
 - Manual: LaTeX (source), PDF & HTML
 - Slides: PowerPoint, PDF and N-up PDF
- Library contains all GSPs discussed in this paper:
 - Mutex[Op]
 - RW[ReadOp, WriteOp]
 - ProdCons[PutOp, GetOp, OtherOp]
 - BoundedProdCons(int size)[PutOp, GetOp, OtherOp]
- GSPs are implemented for multiple thread packages:
 - POSIX, Solaris, Windows, DCE
 - Dummy (for non-threaded applications)

Using GSP classes

- Define one of the following preprocessor symbols before you `#include` a GSP header file
 - `P_USE_POSIX_THREADS`
 - `P_USE_SOLARIS_THREADS`
 - `P_USE_WIN32_THREADS`
 - `P_USE_DCE_THREADS`
 - `P_USE_NO_THREADS`
- Typically done with `-D<symbol>` command-line option to compiler

Summary

- GSPs are a generalization of the `ScopedMutex` class:
 - Out-of-the-box support for mutual-exclusion, readers-writer and (bounded) producer-consumer policies
 - You can write new GSPs if the need arises
- Benefits:
 - Makes it trivial to add synchronization to a C++ class
 - Makes code easier to read and maintain
 - Portability across multiple thread packages
 - Minimal performance overhead due to `inline` implementation
- All software and documentation is available:
 - MIT-style license (open-source, non-viral)
 - Download from www.CiaranMcHale.com/download